Synchronization Requirements of Token Smart Contracts

Giorgia Azzurra Marson NEC Labs Europe

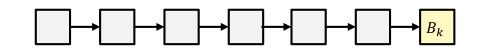
Based on joint work with: Orestis Alpos, Christian Cachin, Luca Zanolini University of Bern

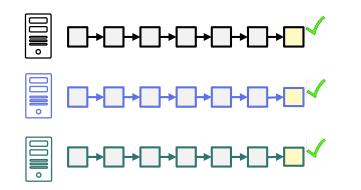
> DLT Workshop 2023 Bologna, Italy

Motivation

Decentralized applications rely on a distributed protocol emulating a shared ledger (blockchain)

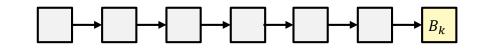
Motivation

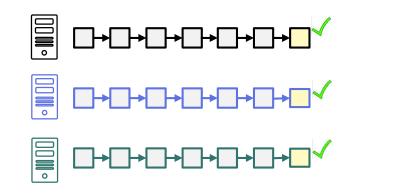


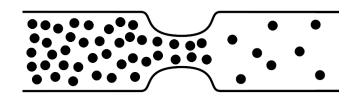


- Decentralized applications rely on a distributed protocol emulating a shared ledger (blockchain)
- Distributed consensus (a.k.a. total-order broadcast) ensures consistency among ledgers

Motivation







transactions submitted by users

transactions confirmed by the blockchain

- Decentralized applications rely on a distributed protocol emulating a shared ledger (blockchain)
- Distributed consensus (a.k.a. total-order broadcast) ensures consistency among ledgers
- However, consensus is the bottleneck of blockchain speed ⊗

Prior work [GKMPS'19]

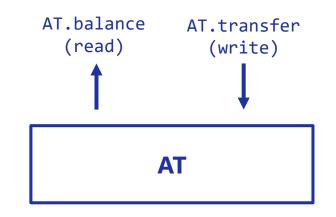
[GKMPS'19] R. Guerraoui, P. Kuznetsov, M. Monti, M. Pavlovic, D.A. Seredinschi: The consensus number of a cryptocurrency. PODC 2019

Prior work [GKMPS'19]

Approach:

- Define AT abstraction as shared-memory object
- Analyze the synchronization power (consensus number) of AT

<u>Main result:</u> The consensus number of **AT** is 1 (range: [1,∞]). "Basic cryptocurrency functionality" = Asset Transfer (AT) object



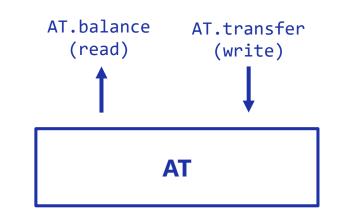
[GKMPS'19] R. Guerraoui, P. Kuznetsov, M. Monti, M. Pavlovic, D.A. Seredinschi: The consensus number of a cryptocurrency. PODC 2019

Prior work [GKMPS'19]

Approach:

- Define AT abstraction as shared-memory object
- Analyze the synchronization power (consensus number) of AT

<u>Main result:</u> The consensus number of **AT** is 1 (range: [1, ∞]). "Basic cryptocurrency functionality" = Asset Transfer (AT) object



Interpretation:

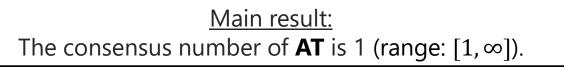
- AT has weakest synchronization power
- Consensus is an overkill for basic cryptocurrency applications
- Intuitive reason: total order is not necessary to prevent double spending (causal order is enough)

[GKMPS'19] R. Guerraoui, P. Kuznetsov, M. Monti, M. Pavlovic, D.A. Seredinschi: The consensus number of a cryptocurrency. PODC 2019

<u>Approach:</u>

- Define AT abstraction as shared-memory object
- Analyze the synchronization power (consensus number) of AT

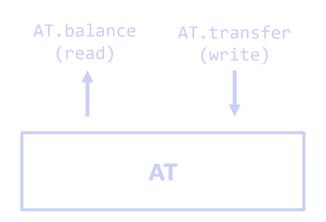
"Basic cryptocurrency functionality" = Asset Transfer (AT) object



What about Smart Contracts?

Interpretation:

- AT has weakest synchronization power
- Consensus is an overkill for basic cryptocurrency applications
- Intuitive reason: total order is not necessary to prevent double spending (causal order is enough)



This work [ACMZ'21]

Motivating question:

What level of synchronization is required for "popular" smart contracts?

This work [ACMZ'21]

<u>Motivating question:</u> What level of synchronization is required for "popular" smart contracts?

ERC-20 token standard

(most popular fungible token in Ethereum)

Approach:

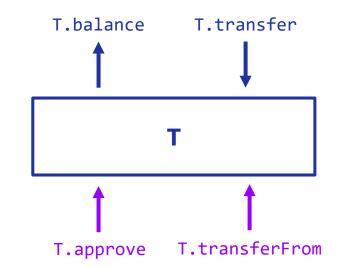
- Define ERC-20 abstraction as shared object T
- Analyze the consensus number of ERC-20 object

This work [ACMZ'21]

Motivating question: What level of synchronization is required for "popular" smart contracts?

ERC-20 token standard

(most popular fungible token in Ethereum)



Approach:

- Define ERC-20 abstraction as shared object T
- Analyze the consensus number of ERC-20 object

New features compared to AT:

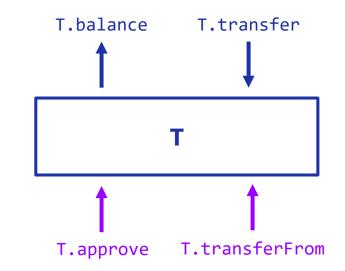
- Account owners can delegate approved spenders to manage asset
- Approval of spenders is dynamic (any time, arbitrary amounts)

This work [ACMZ'21]

Motivating question: What level of synchronization is required for "popular" smart contracts?

ERC-20 token standard

(most popular fungible token in Ethereum)



Approach:

- Define ERC-20 abstraction as shared object T
- Analyze the consensus number of ERC-20 object

New features compared to AT:

- Account owners can delegate approved spenders to manage asset
- Approval of spenders is dynamic (any time, arbitrary amounts)

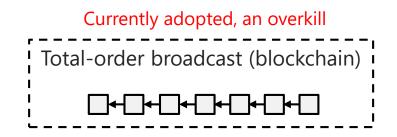
Main results:The consensus number of **ERC-20 token T** dynamically
changes with the contract state (q): $CN(T_q) = 1 + max \{\# \text{ approved spenders for account } a\}$

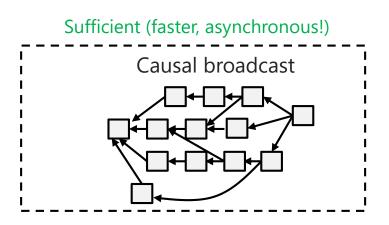
Outlook

<u>Prior work</u> Synchronization power of **cryptocurrency**:

CN(AT) = 1

 \Rightarrow transactions can be processed concurrently \Rightarrow total order is not necessary, causal broadcast can be used instead



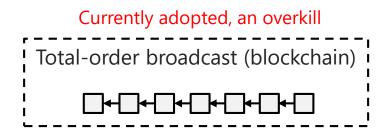


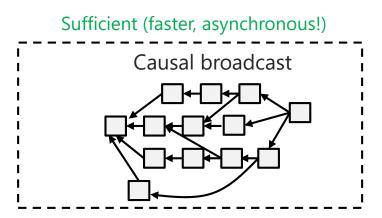
Outlook

<u>Prior work</u> Synchronization power of **cryptocurrency**:

CN(AT) = 1

 \Rightarrow transactions can be processed concurrently \Rightarrow total order is not necessary, causal broadcast can be used instead





<u>This work</u> Synchronization power of Ethereum **ERC-20 token T**:

 $CN(T_q) = 1 + \max_a \{ \# \text{ approved spenders for account } a \}$

⇒ transactions can be processed concurrently, if
 issued by spenders of different accounts
 ⇒ total order is needed only for resolving conflicts,
 causal broadcast could be used optimistically

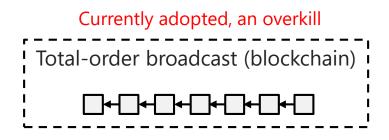
Ideally: optimally-concurrent protocol for "useful" smart contracts As-concurrent-as-possible broadcast

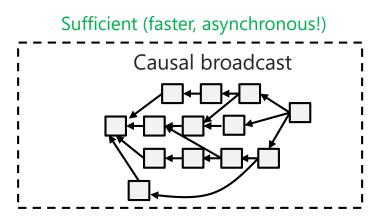
Outlook

<u>Prior work</u> Synchronization power of **cryptocurrency**:

CN(AT) = 1

 \Rightarrow transactions can be processed concurrently \Rightarrow total order is not necessary, causal broadcast can be used instead

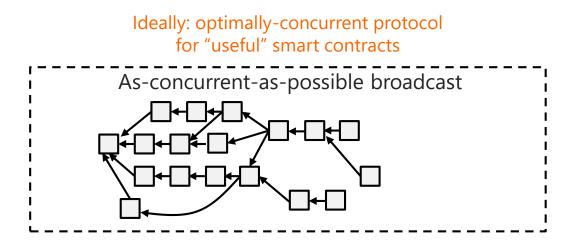




<u>This work</u> Synchronization power of Ethereum **ERC-20 token T**:

 $CN(T_q) = 1 + \max_a \{ \# \text{ approved spenders for account } a \}$

⇒ transactions can be processed concurrently, if
 issued by spenders of different accounts
 ⇒ total order is needed only for resolving conflicts,
 causal broadcast could be used optimistically



Thank you for your attention ©

Backup slide



Synchronization power (consensus number) of shared objects

Wait-free hierarchy [Herlihy'91]

- Consensus is universal: any shared object has a wait-free implementation from consensus objects
- ⇒ consensus can serve as reference for the synchronization power of shared objects

Consensus number of object O: CN(O) :=max $n \mid \exists$ wait-free implementation of consensus object from objects of type O and registers, in a system with n processes.

- Intuitively: max # processes that can be synchronized "using O"
 CN(O) = 1 ⇒ O useless for synchronization
 CN(O) = ∞ ⇒ O can synchronize any number of processes
- Metric to compare synchronization power of shared objects

[Herlihy'91] M. Herlihy: Wait-Free Synchronization. ACM Trans. Program. Lang. Syst. 1991

